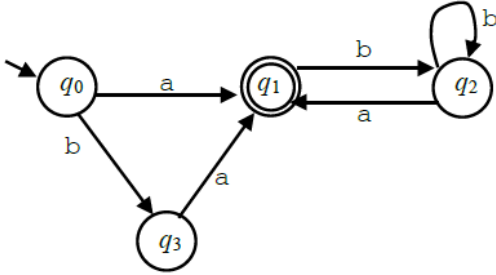
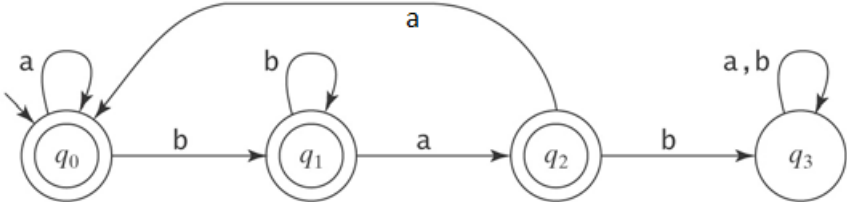
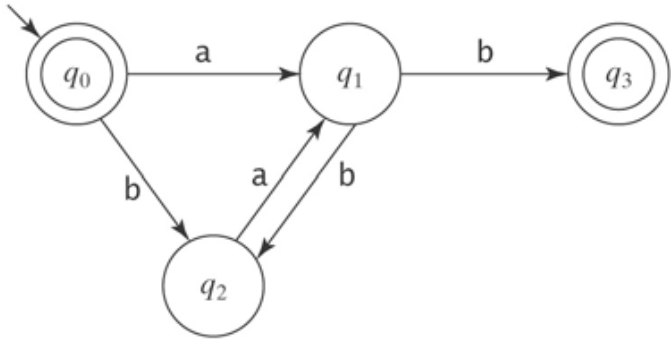
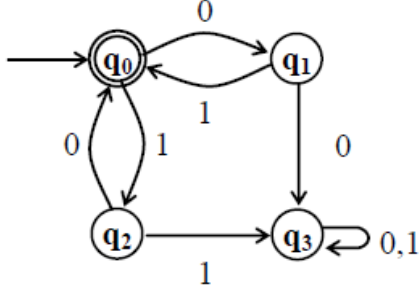
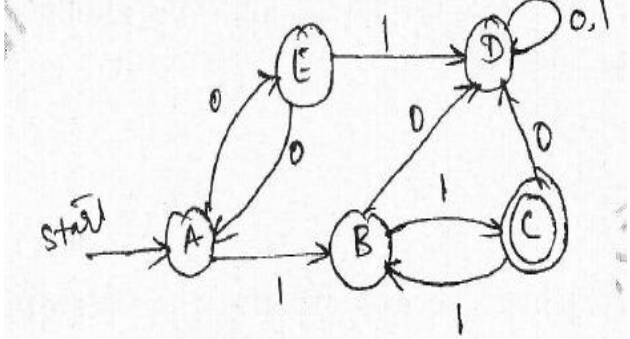
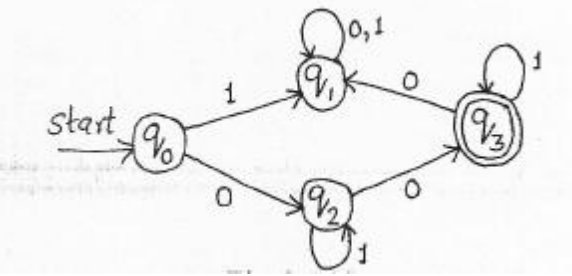
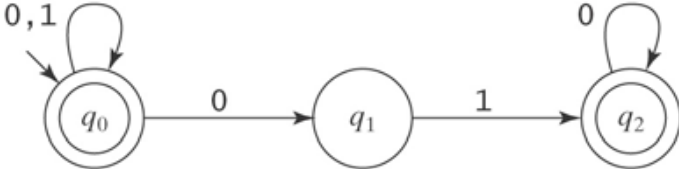


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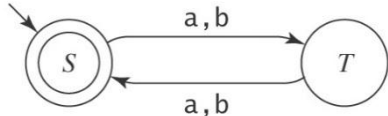
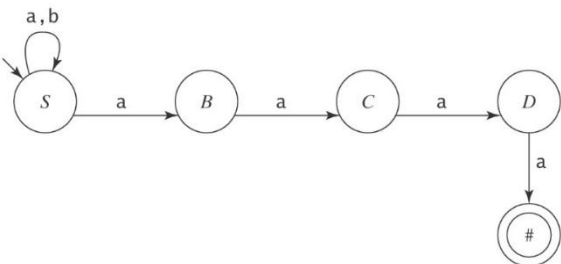
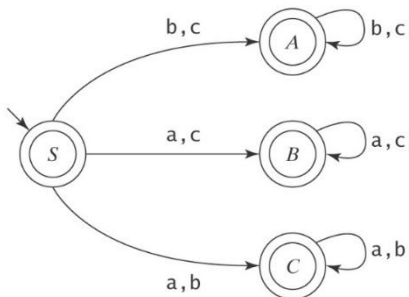
Module-2: Review Questions (up to IA1)

Sl. No.	Questions	Marks
Regular Expression		
1.	Define regular expression.	5
2.	<p>Write the RE for following languages.</p> <p>a. $\{w \in \{a, b\}^* : \text{every } a \text{ in } w \text{ is immediately preceded and followed by } b\}$.</p> <p>b. $\{w \in \{a, b\}^* : w \text{ does not end in } ba\}$.</p> <p>c. $\{w \in \{0, 1\}^* : w \text{ corresponds to the binary encoding, without leading } 0\text{'s, of natural numbers that are evenly divisible by } 4\}$.</p> <p>d. $\{w \in \{0, 1\}^* : w \text{ corresponds to the binary encoding, without leading } 0\text{'s, of natural numbers that are powers of } 4\}$.</p> <p>e. $\{w \in \{0-9\}^* : w \text{ corresponds to the decimal encoding, without leading } 0\text{'s, of an odd natural number}\}$.</p> <p>f. $\{w \in \{0, 1\}^* : w \text{ has } 001 \text{ as a substring}\}$.</p> <p>g. $\{w \in \{a, b\}^* : w \text{ has both } aa \text{ and } bb \text{ as substrings}\}$.</p> <p>h. $\{w \in \{0, 1\}^* : \text{none of the prefixes of } w \text{ ends in } 0\}$.</p> <p>i. $\{w \in \{a, b\}^* : \#_a(w) \equiv_3 0\}$.</p> <p>j. $\{w \in \{a, b\}^* : \#_a(w) \leq 3\}$.</p> <p>k. $\{w \in \{a, b\}^* : w \text{ contains no more than two occurrences of the substring } aa\}$.</p> <p>l. $\{a^{2n}b^{2m} \mid m \geq 0, n \geq 0\}$</p> <p>m. $\{a^n b^m \mid n \geq 4, m \leq 3\}$</p> <p>n. $\{w : w \bmod 3 = 0 \text{ where } w \in (a,b)^*\}$</p> <p>o. strings of 0's and 1's having no two consecutive 0's.</p> <p>p. strings of a's and b's starting with 'a' and ending with 'b'</p>	3 each
3.	Let L be the language accepted by the following finite state machine:	4

	<p>Indicate, for each of the following regular expressions, whether it correctly describes L:</p> <ol style="list-style-type: none"> $(a \cup ba)bb^*a$. $(\epsilon \cup b)a(bb^*a)^*$. $ba \cup ab^*a$. $(a \cup ba)(bb^*a)^*$. 	
4.	State Kleene's theorem.	2
5.	<p>Convert following RE's into FSM's</p> <ol style="list-style-type: none"> $(ab^*)^*$ $(a \cup b)^*$ $a^* \cup b^*$ $(a \cup b)ab$ 	3 each
6.	<p>Convert following FSM's to RE's</p> <p>(a)</p>  <p>(b)</p> 	6 each

	<p>c)</p>  <p>d)</p>  <p>e)</p> 	
7.	<p>Given the following DFSM M, write a regular expression that describes $L(M)$:</p> 	
6.	Briefly explain the applications of regular expression.	6
7.	<p>Simplify each of the following regular expressions:</p> <ol style="list-style-type: none"> $(a \cup b)^* (a \cup \epsilon) b^*$. $(\emptyset^* \cup b) b^*$. $(a \cup b)^* a^* \cup b$. $((a \cup b)^*)^*$. $((a \cup b)^+)^*$. $a((a \cup b)(b \cup a))^* \cup a((a \cup b)a)^* \cup a((b \cup a)b)^*$. 	3 each

Regular Grammars

8.	Define regular grammar	2
	<p>Give regular grammar for each of the following languages:</p> <p>a) $\{w \in \{a, b\}^* : w \text{ contains an odd number of } a\text{'s and an odd number of } b\text{'s}\}.$</p> <p>b) $\{w \in \{a, b\}^* : w \text{ does not end in } aa\}.$</p> <p>c) $\{w \in \{a, b\}^* : w \text{ contains the substring } abb\}.$</p> <p>d) $\{w \in \{a, b\}^* : \text{if } w \text{ contains the substring } aa \text{ then } w \text{ is odd}\}.$</p> <p>e) $\{w \in \{a, b\}^* : w \text{ does not contain the substring } aabb\}.$</p> <p>f) $\{w \in \{a, b\}^* : \text{every } a \text{ in } w \text{ is immediately followed by at least one } b\}.$</p>	3 or 4 each
9.	<p>Give regular grammar for following FSM's</p> <p>(a)</p>  <p>(b)</p>  <p>(c)</p> 	6 each
10.	<p>Construct FSM for the following regular grammar</p> <p>(a) $S \rightarrow aT$ $T \rightarrow bT$ $T \rightarrow aT \rightarrow aW$ $W \rightarrow \epsilon$ $W \rightarrow aT$</p> <p>(b) $S \rightarrow bS$ $S \rightarrow aT$ $S \rightarrow \epsilon$ $T \rightarrow bS$</p>	4 each

Regular and Non-Regular Languages (Pumping Theorem)		
11.	State and prove pumping theorem for regular languages	6
12.	List the applications of pumping lemma	4
13.	Prove that “The regular languages are closed under complement, intersection, difference, reverse, and letter substitution”	6
14.	Prove that following languages are not regular using pumping theorem <ul style="list-style-type: none"> • $L = \{a^n b^n \mid n \geq 0\}$ • Balanced parenthesis language • Even palindrome language (palindrome over {a,b} with even length) • The language with more a's than b's • The prime number of a's, $L = a^n$ where n is prime 	8 each